

Hardware Benchmarking of Round 2 Candidates in the NIST Lightweight Cryptography Standardization Process

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NIST Standardization Process

- Performance of current NIST cryptographic standards not acceptable in constrained environments (e.g., sensor networks, healthcare, the Internet of Things, cyber physical systems)

Timeline of the NIST Lightweight Cryptography Standardization:

- Aug. 2018: Submission Requirements and Evaluation Criteria
- Feb. 2019: 57 candidates submitted
- Aug. 2019: 32 candidates qualified for Round 2
- Feb. 2021: Decision on Round 3 candidates expected

Hardware Benchmarking Goals

- Stimulate the development of hardware implementations that can be fairly compared with each other (e.g., common API & development package)
- Perform design space exploration of at least selected candidates
- Evaluate and rank candidates from the point of view of their performance in hardware

Benchmarking Platforms

- Widely used low-cost, low-power FPGA families
- Capable of holding side-channel-protected designs (possibly using up to 4 times more resources than unprotected designs)
- Supported by free versions of state-of-the-art industry tools



- Xilinx: **Artix-7** : xc7a12tcsg325-3 (smallest)
- Intel: **Cyclone 10 LP** : 10CL016-YF484C6
- Lattice Semiconductor: **ECP5** : LFE5U-25F-6BG381C

Optimization Target

- **Maximum Throughput assuming**
 - **Up to 2500 LUTs, 5000 flip-flops of Artix-7 FPGA**
 - **No BRAMs & no DSP units**
 - **Resources comparable to those used by the lightweight implementation of the current standard AES-GCM**

Benchmarking Metrics

1. Resource Utilization

- Number of LUTs (LEs for Cyclone 10LP)

2. Throughput in Mbits/s

- for the following sizes of inputs
 - a. long [with Throughput = $d \cdot \text{Block size} / (\text{Time}(N+d \text{ blocks}) - \text{Time}(N \text{ blocks}))$]
 - b. 1536 bytes
 - c. 64 bytes
 - d. 16 bytes.
- all throughputs calculated separately for
 - authenticated encryption: Plaintext, Associated Data (AD)
 - hashing

Summary of Hardware Design Submissions

32 submissions representing 25 out of 32 candidates (78%)

Candidate with 3 independent submissions:

Xoodyak

Candidates with 2 independent submissions:

Ascon, COMET, Gimli, Subterranean 2.0, TinyJAMBU

8 submissions from George Mason University

24 by groups from all over the world

Design Variants

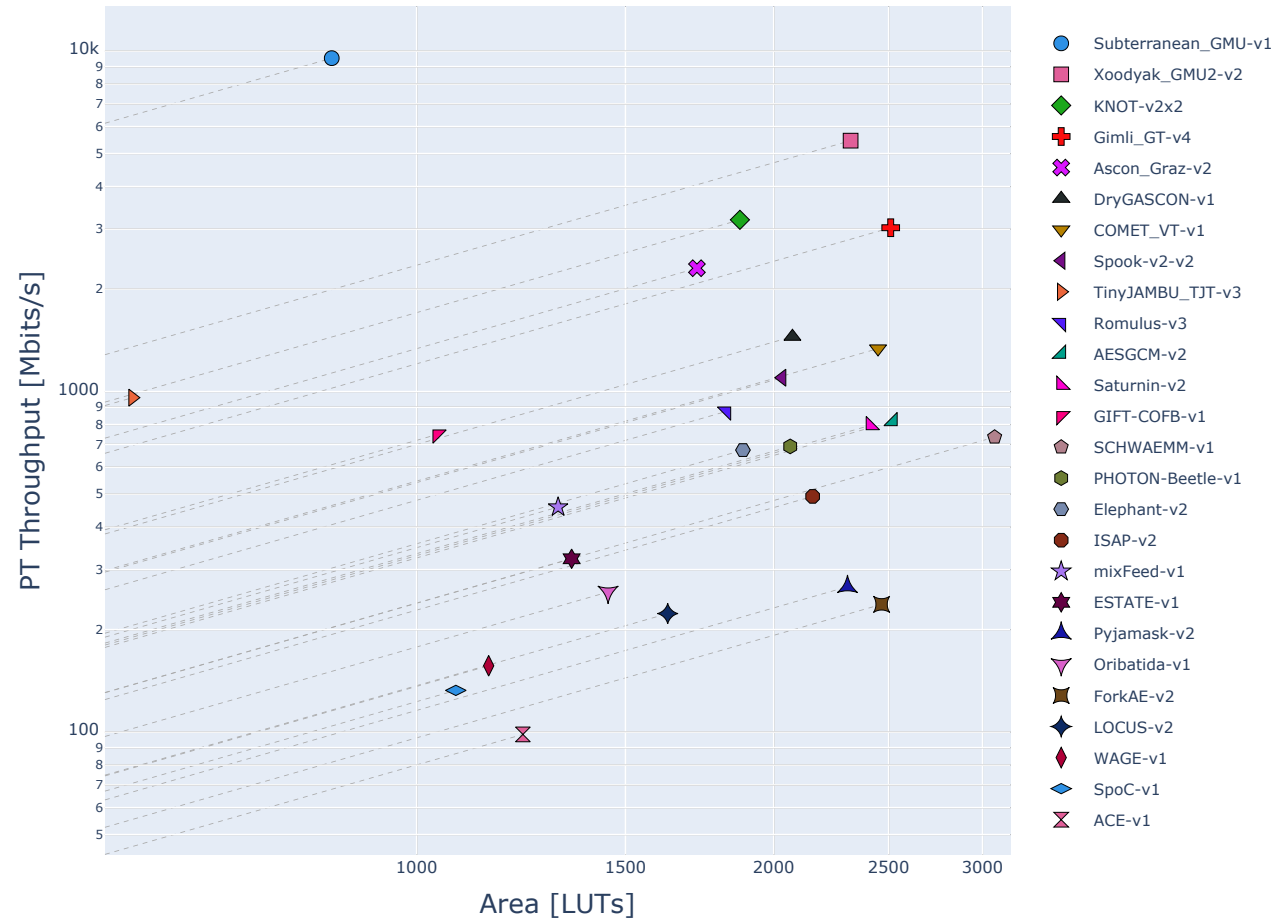
Different variants correspond to

- different algorithms of the same family
- different parameter sets, such as sizes of keys, nonces, tags, etc.
- support for authenticated encryption vs. authenticated encryption+hashing
- different hardware architectures, e.g., basic iterative, folded, unrolled

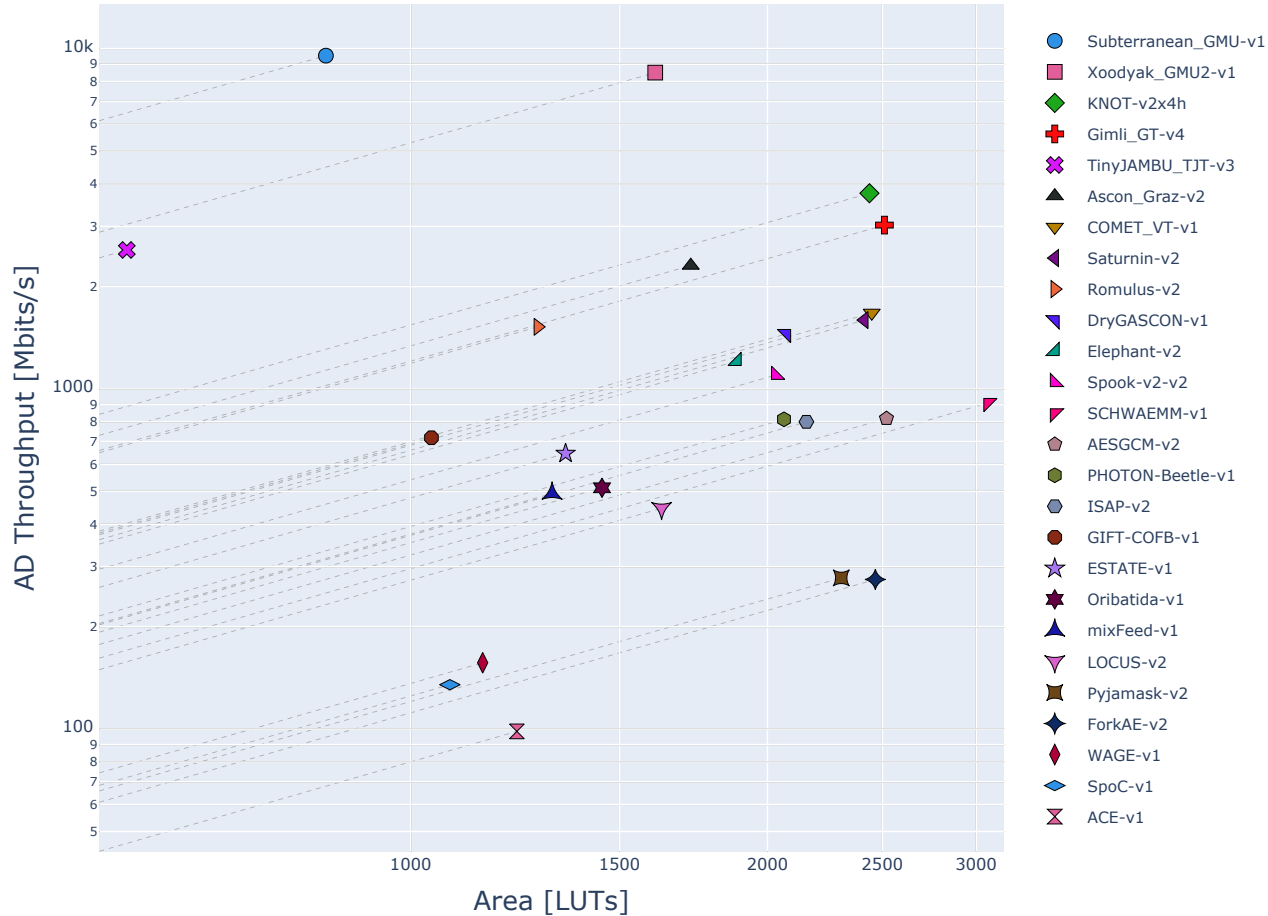
92 variants

Minimum: 1, Maximum: 16, Average: 3.1
per hardware design submission

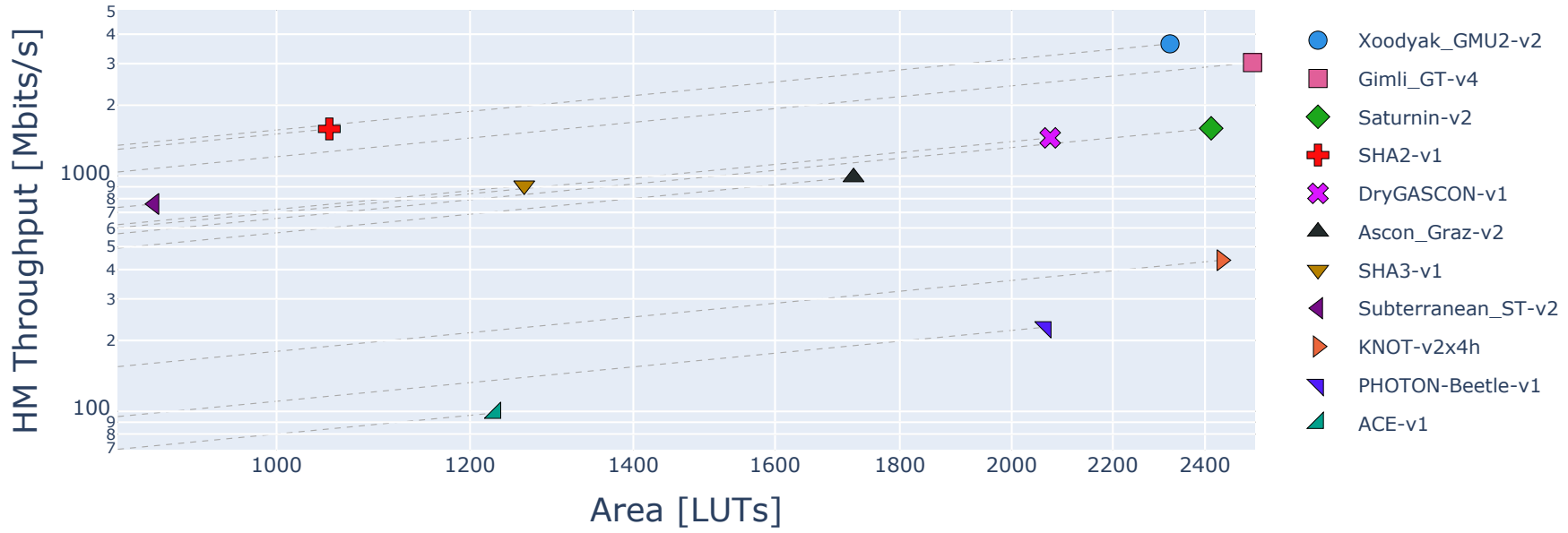
Throughput vs. Area for Long Plaintext: Artix-7



Throughput vs. Area for Long AD: Artix-7



Throughput vs. Area for Hashing: Artix-7



Dependence of Ranking on Input Size

| Position | Long | 1536 B | 64 B | 16 B |
|----------|--------------|--------------|--------------|---------------|
| 1 | Subterranean | Subterranean | Subterranean | Subterranean |
| 2 | Xoodyak | Xoodyak | Xoodyak | Xoodyak |
| 3 | KNOT | KNOT | KNOT | Ascon |
| 4 | Gimli | Gimli | Ascon | COMET |
| 5 | Ascon | Ascon | DryGASCON | DryGASCON |
| 6 | DryGASCON | DryGASCON | Gimli | KNOT |
| 7 | COMET | COMET | COMET | TinyJAMBU |
| 8 | Spook v2 | Spook v2 | TinyJAMBU | Romulus |
| 9 | TinyJAMBU | TinyJAMBU | Romulus | Gimli |
| 10 | Romulus | Romulus | Spook v2 | PHOTON-Beetle |

Higher position
Lower position
for smaller
messages

Conclusions

- For **authenticated encryption of plaintexts** 10 candidates outperform AES-GCM: **Subterranean 2.0, Xoodyak, KNOT, Gimli, Ascon, DryGASCON, COMET, Spook v2, TinyJAMBU, and Romulus.**
- For **processing of associated data (ADs)** all of them, as well as **Saturnin and Elephant**, outperform AES-GCM
- Out of them:
 - **Xoodyak, Gimli, and Saturnin** support hashing faster than SHA-2
 - **DryGASCON and Ascon**, perform hashing faster than the folded implementation of SHA-3
- All of the mentioned above 12 candidates have good chances of qualifying for Round 3

Concurrent & Future Work

- Evaluation in terms of Power consumption and Energy per bit
- ASIC Benchmarking
- Side-channel protected implementations of Round 3 candidates

Most recent results:

Cryptology ePrint Archive: Report 2020/1207

<https://cryptography.gmu.edu/athena>